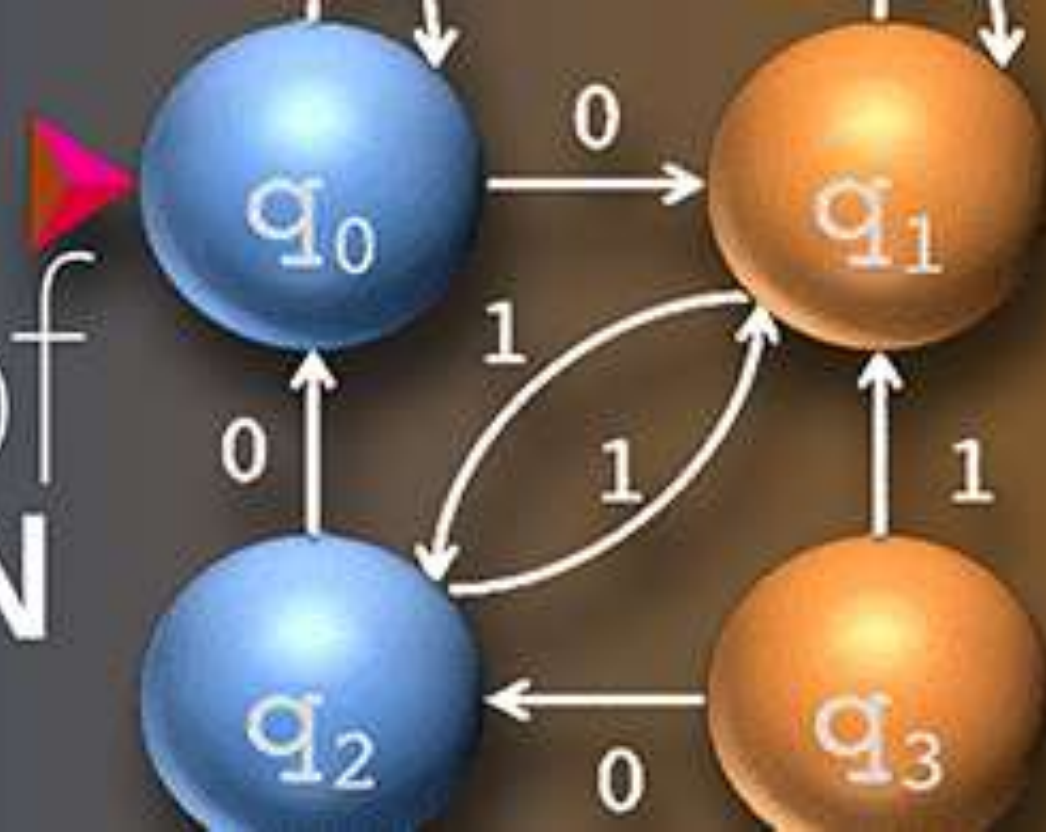


CSE 305

Theory of COMPUTATION



Lecture 9

Mathematical Preliminaries (5)



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Contents

Mathematical Preliminaries



- Set Theory
- Sequences and Tuples
- Relations and Functions
- Alphabets, Strings and Languages**
- Graph and Tree**
- Mathematical Logic**

Alphabets, Strings and Languages

- **Alphabets:** The alphabet of any language is defined as a finite, non-empty set of symbols. These form the basic units of strings. Alphabets may be denoted by Σ .
 - Some examples of alphabets include: binary alphabets, $\Sigma = \{0, 1\}$; English alphabets, $\Sigma = \{A, B, C, \dots, Z\}$, numbers, $\Sigma = \{0, 1, 2, \dots, 9\}$, etc.
- **Strings:** The symbols from an alphabet can be combined to form words or strings. A string is defined as a finite sequence of symbols chosen from some alphabets over Σ .
 - For example: (i) 01, 011, 0000, 101011 are all strings over an alphabet $\Sigma = \{0, 1\}$; (ii) aa, ac, abc, aabb are all strings over an alphabet $\Sigma = \{a, b, c\}$; (iii) 10, 123, 1024 are all strings over an alphabet $\Sigma = \{0, 1, 2, 3, \dots, 9\}$.

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Alphabets, Strings and Languages

Languages:

- The finite or infinite set of all strings obtained from the alphabet Σ is called **language**. A language is a subset of Σ^* , and defined as-

$$L = \{w \in \Sigma^* \mid w \text{ is a string over } \Sigma \text{ and has some specific property}\}$$

- **Example:**

(i) $L_1 = \{0, 00, 101, 1100\}$ is a language over the alphabet $\Sigma = \{0, 1\}$;

(ii) $L_2 = \{ww^R \mid w \in \{a, b\}^*\}$ is a language over the alphabet $\Sigma = \{a, b\}$.

Alphabets, Strings and Languages

Operations on Languages:

- **Union of Languages:** The union of languages is similar to the union of sets. The union of two languages is defined as-

$$L_1 \cup L_2 = \{x \mid x \in L_1 \text{ or } x \in L_2\}$$

- **Intersection of Languages:** The intersection of languages is similar to the intersection of sets. The intersection of two languages is defined as-

$$L_1 \cap L_2 = \{x \mid x \in L_1 \text{ and } x \in L_2\}$$

- **Concatenation of Languages:** Concatenation of languages involves grouping all possible strings formed by combining strings of two different languages. Concatenation of two languages is defined as-

$$L_1 L_2 = \{xy \mid x \in L_1 \text{ and } y \in L_2\}$$

Alphabets, Strings and Languages

Operations on Languages:

- **Reversal of Languages:** The reversal of a language L is denoted by L^R and it is defined as-

$$L^R = \{x \mid x^R \in L\}$$

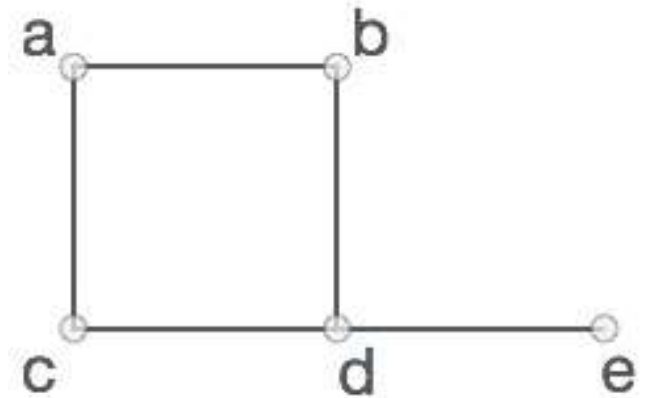
Reversal of a language L is defined as a collection of reversal of all strings of L .

- **Kleene Star (Kleene Closure):** The Kleene star of a language L is defined as the language L^* consisting of all the strings obtained by concatenating any finite number (including zero) of strings from L together.

For example, the Kleene star of a language L given by $\{0, 1\}$ may be written as $L^* = \{\epsilon, 0, 1, 10, 00, 01, 11, 100, 110, 101, \dots\}$

Graphs

- A graph is a pictorial representation of a set of objects where some pairs of objects are connected by links. The interconnected objects are represented by points termed as **vertices**, and the links that connect the vertices are called **edges**.
- Formally, a graph G is a pair of sets (V, E) , where V is the set of vertices and E is the set of edges, connecting the pairs of vertices, such as $G = (V, E)$:
 1. A set V of elements called nodes (or points or vertices)
 2. A set E of edges, such that each edge e in E is identified with a unique (unordered) pair $[u, v]$ of nodes in V , denoted by $e = [u, v]$; the nodes u and v are called endpoints of e , and u and v are said to be adjacent nodes or neighbors.



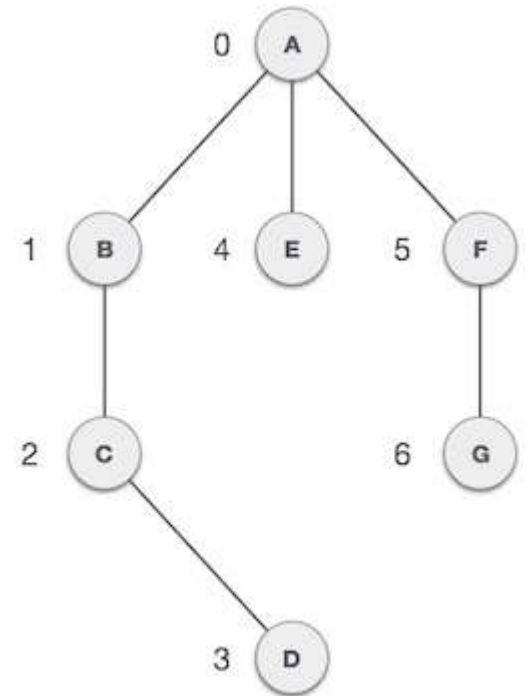
- In the given graph:
 $V = \{a, b, c, d, e\}$
 $E = \{ab, ac, bd, cd, de\}$



Graphs

Graph Terminology:

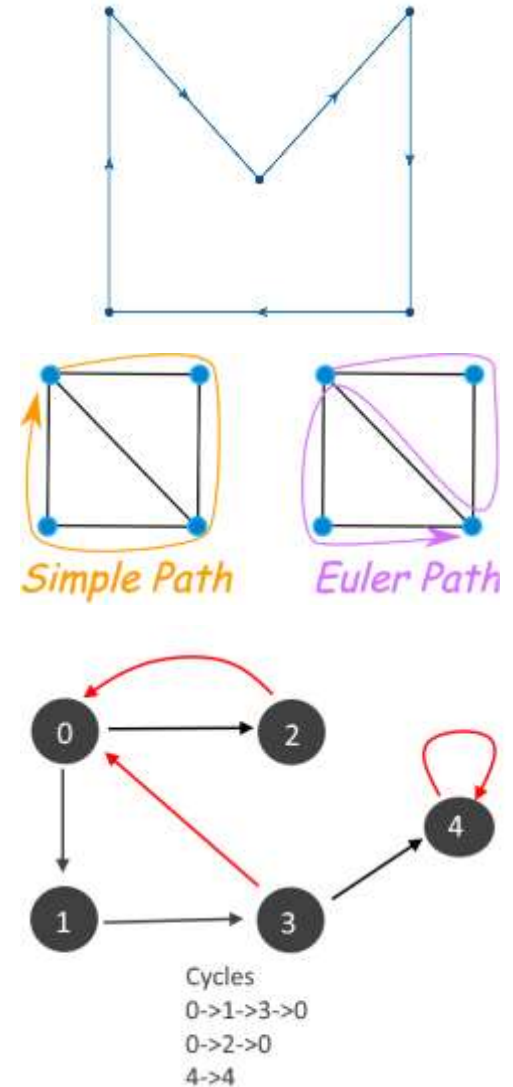
- **Vertex** – Each node of the graph is represented as a vertex. In the example, the labeled circle represents vertices. Thus, A to G are vertices.
- **Edge** – Edge represents a path between two vertices or a line between two vertices. In the example, the lines from A to B, B to C, and so on represents edges.
- **Adjacency** – Two node or vertices are adjacent if they are connected to each other through an edge. In the example, B is adjacent to A, C is adjacent to B, and so on.
- **Path** – Path represents a sequence of edges between the two vertices. In the example, ABCD represents a path from A to D.



Graphs

Graph Terminology:

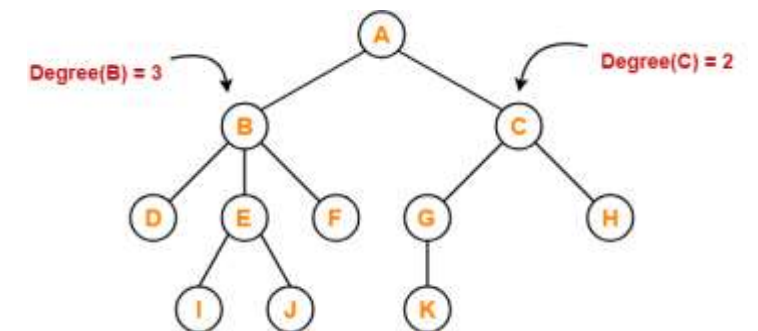
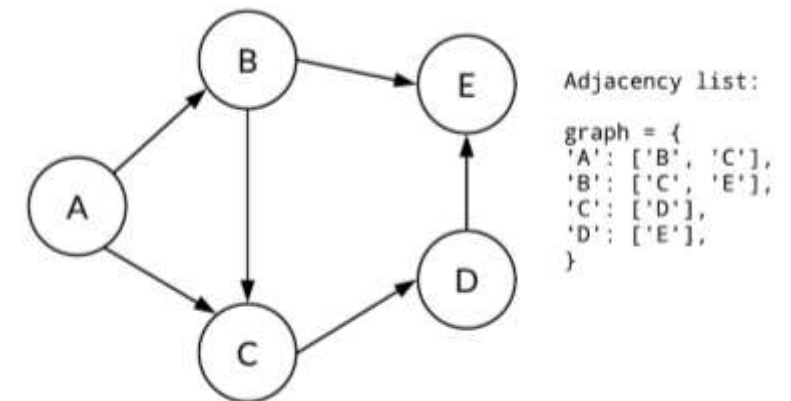
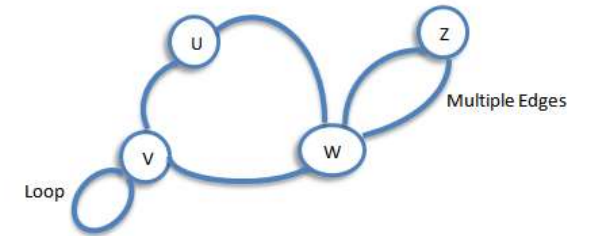
- **Closed Path** - A path will be called as closed path if the initial node is same as terminal node. A path will be closed path if $V_0=V_N$.
- **Simple Path and Euler Path** - If all the nodes of the graph are distinct with an exception $V_0=V_N$, then such path P is called as closed simple path. A path is called an **Euler path** if it traverses each edge lying in it exactly once. A path is called an **Euler circuit** if it traverses each edge lying in it exactly once and whose initial and final vertices are the same.
- **Cycle** - A cycle can be defined as the path which has no repeated edges or vertices except the first and last vertices.



Graphs

Graph Terminology:

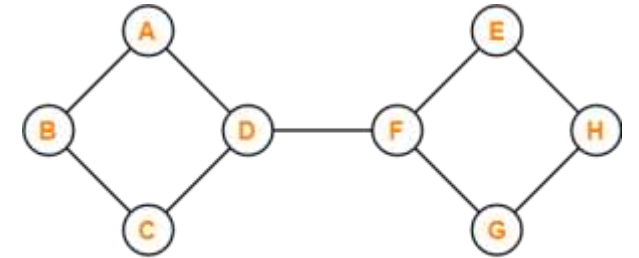
- **Loop** - An edge that is associated with the similar end points can be called as Loop.
- **Adjacent Nodes** - If two nodes u and v are connected via an edge e , then the nodes u and v are called as neighbours or adjacent nodes.
- **Degree of the Node** - A degree of a node is the number of edges that are connected with that node. A node with degree 0 is called as isolated node.



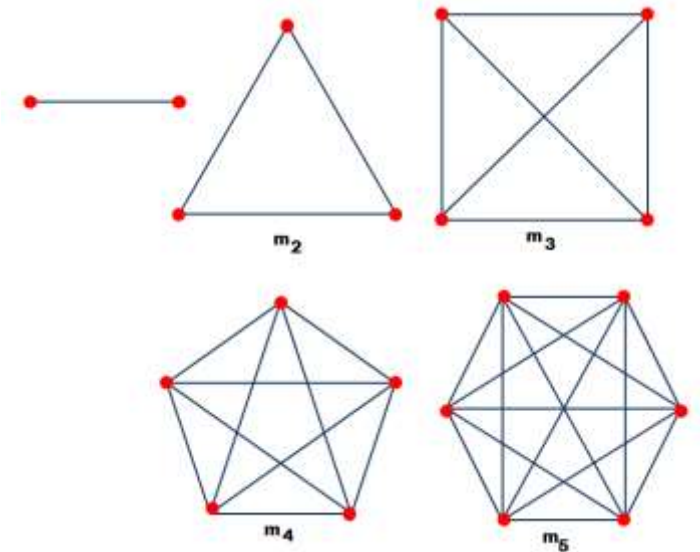
Graphs

Types of Graphs:

- **Connected Graph** - A connected graph is the one in which some path exists between every two vertices (u, v) in V . There are no isolated nodes in connected graph.
- **Complete Graph** - A complete graph is the one in which every node is connected with all other nodes. A complete graph contain $n(n-1)/2$ edges where n is the number of nodes in the graph.



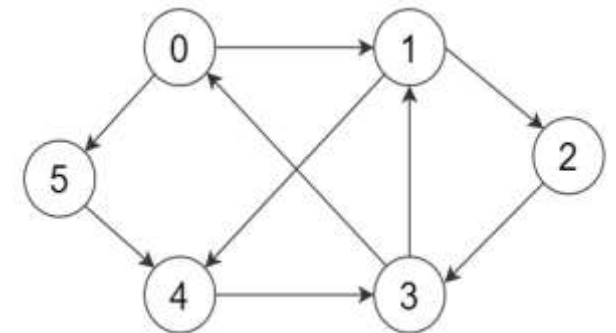
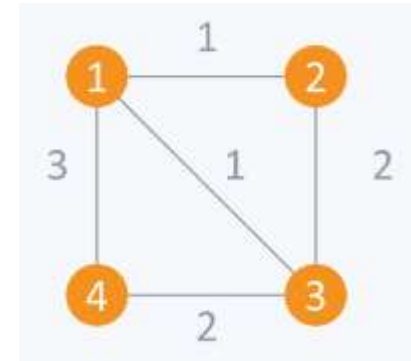
Example of Connected Graph



Graphs

Types of Graphs:

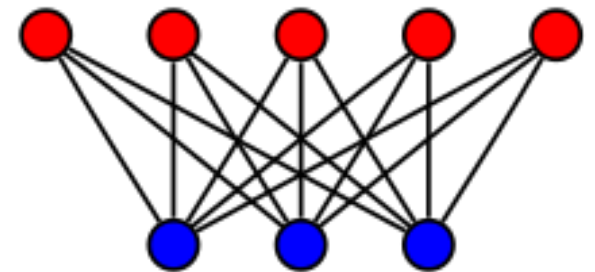
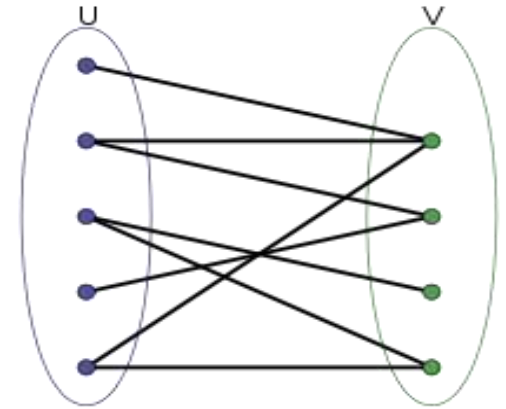
- **Weighted Graph** - In a weighted graph, each edge is assigned with some data such as length or weight. The weight of an edge e can be given as $w(e)$ which must be a positive (+) value indicating the cost of traversing the edge.
- **Digraph** - A digraph is a directed graph in which each edge of the graph is associated with some direction and the traversing can be done only in the specified direction.



Graphs

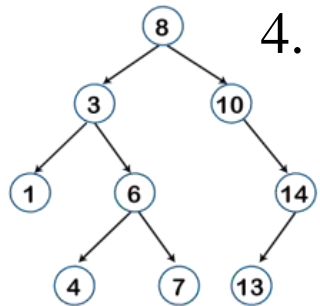
Types of Graphs:

- **Bipartite Graph** - A graph $G = (V, E)$ is called a bipartite graph (also known as a biograph) if its vertices can be decomposed into two disjoint sets such that no two graph vertices within the same set are adjacent. A bipartite graph is a special case of a k -partite graph with .
- **Complete Bipartite Graph** - A graph $G = (V, E)$ is called a complete bipartite graph if its vertices V can be partitioned into two subsets V_1 and V_2 such that each vertex of V_1 is connected to each vertex of V_2 . The number of edges in a complete bipartite graph is $m*n$, as each of the m vertices is connected to each of the n vertices.

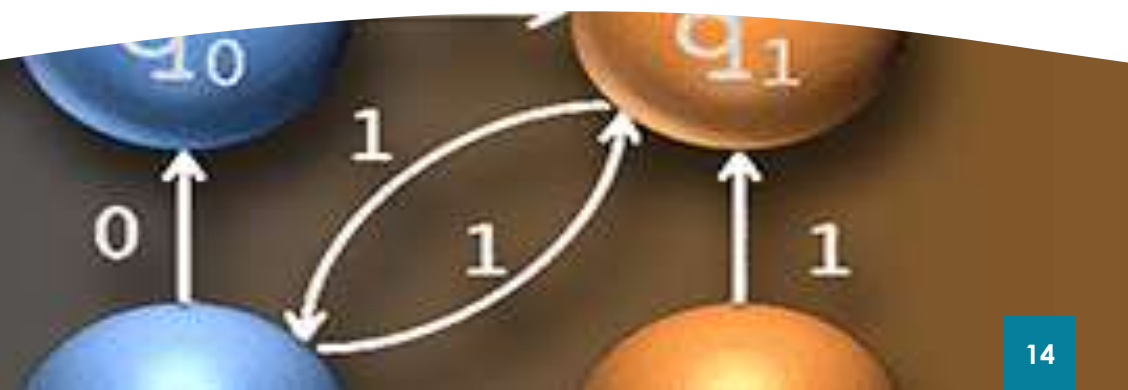


Trees

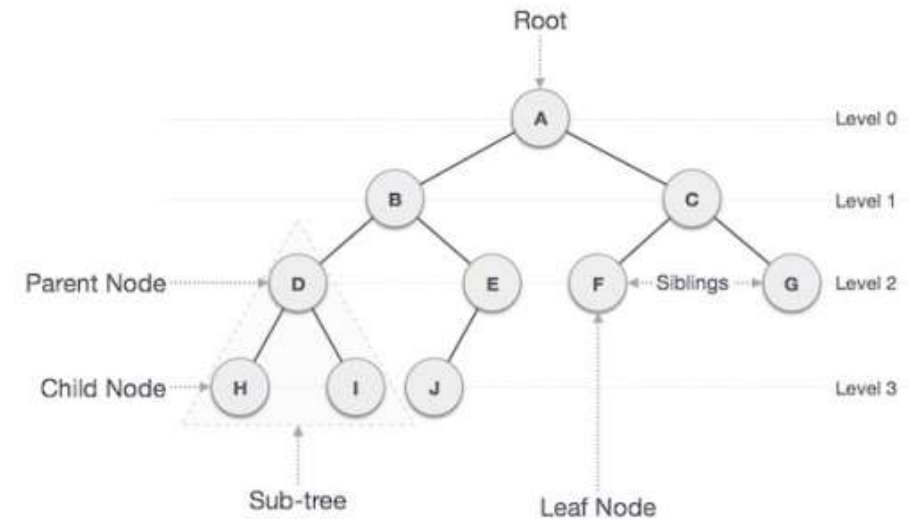
- **Tree** is mainly used to represent data containing a hierarchical relationship between nodes/elements, connected by **edges**. Tree is a special kind of graph that does not contain any cycle/loop. A tree can also be defined as a non-empty, finite set of elements/nodes, which posses the following properties:
 1. There is a special node called the **root** node of the tree. The root has no incoming edges.
 2. The remaining nodes of the tree form an ordered pair of disjoint sub-trees, if it is a binary tree.
 3. There is exactly one path from the root to every other node in the tree.
 4. The nodes that do not have any outgoing edges are called **leaves** of the tree.



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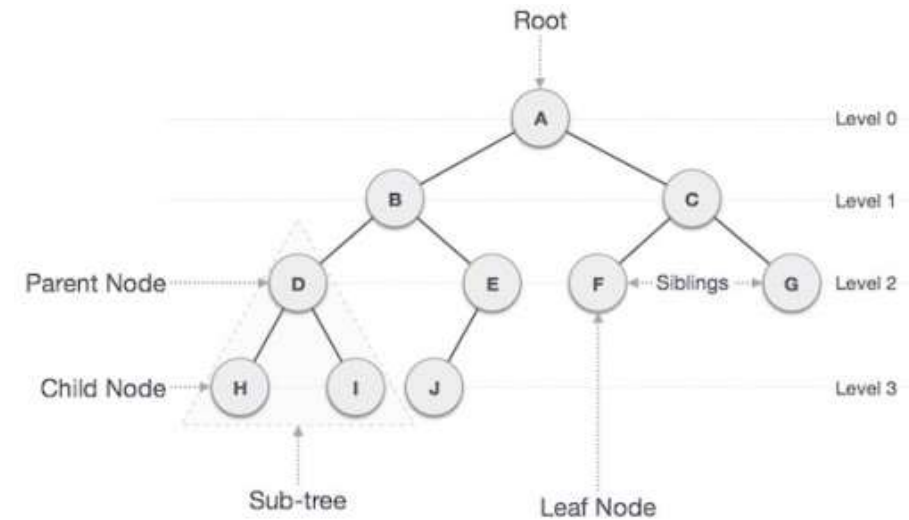
Trees



Tree Terminology:

- **Path** – Path refers to the sequence of nodes along the edges of a tree.
- **Root** – The node at the top of the tree is called root. There is only one root per tree and one path from the root node to any node.
- **Parent** – Any node except the root node has one edge upward to a node called parent.
- **Child** – The node below a given node connected by its edge downward is called its child node.
- **Sibling** - The nodes that have the same parent are known as siblings.
- **Leaf** – The node which does not have any child node is called the leaf node.

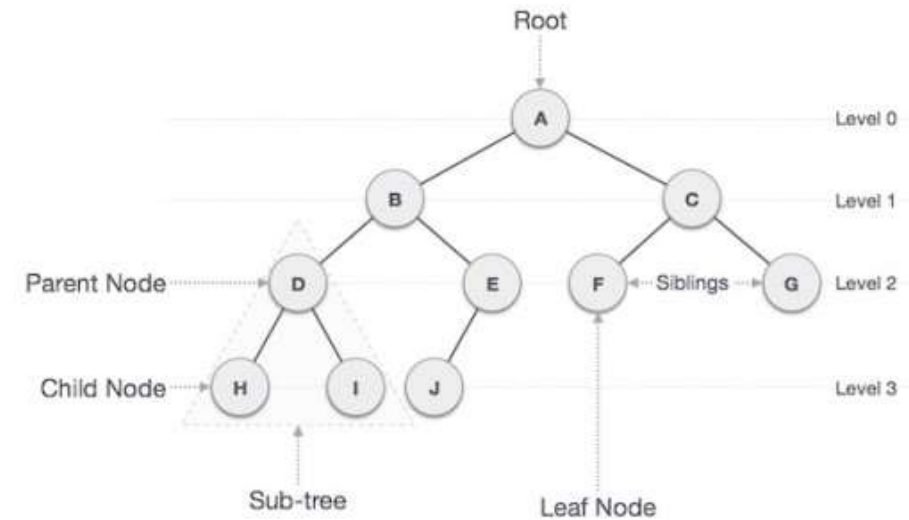
Trees



Tree Terminology:

- **Internal nodes:** A node has at least one child node known as an *internal*
- **Subtree** – Subtree represents the descendants of a node.
- **Ancestor node:** An ancestor of a node is any predecessor node on a path from the root to that node. The root node doesn't have any ancestors.
- **Descendant:** The immediate successor of the given node is known as a descendant of a node.
- **Visiting:** Visiting refers to checking the value of a node when control is on the node.

Trees



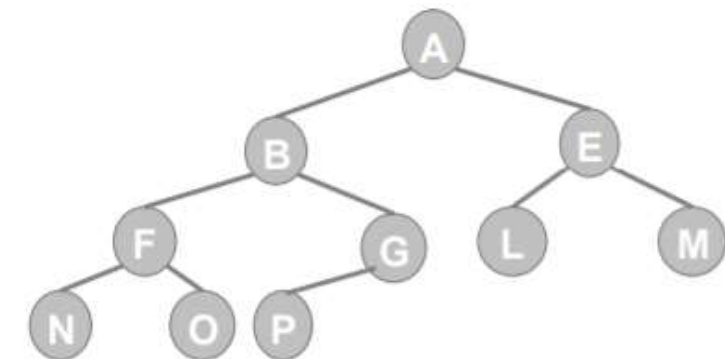
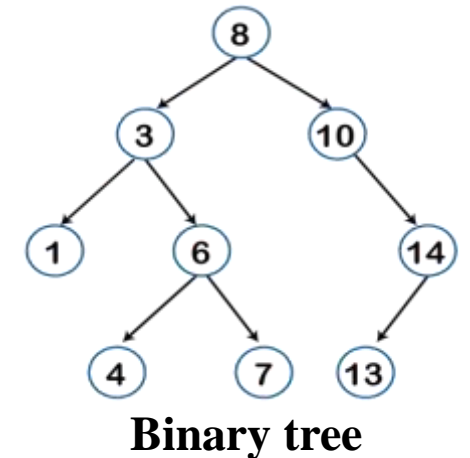
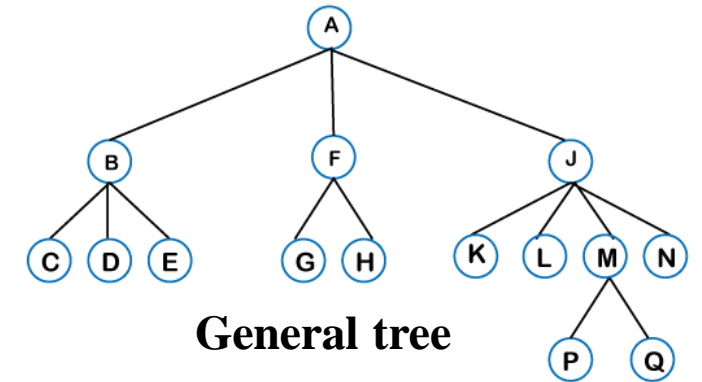
Tree Terminology:

- **Traversing** – Traversing means passing through nodes in a specific order.
- **keys** – Key represents a value of a node based on which a search operation is to be carried out for a node.
- **Levels** – Level of a node represents the generation of a node. If the root node is at level 0, then its next child node is at level 1, its grandchild is at level 2, and so on.

Trees

Types of Trees:

- **General Tree:** A tree where a node can have any number of children/descendants is called a general tree. A node can have either 0 or maximum n number of nodes.
- **Binary Tree:** In a binary tree, each node in a tree can have utmost two child nodes. Here, utmost means whether the node has 0 nodes, 1 node or 2 nodes.
- **Complete Binary Tree:** A binary tree is said to be complete, if all its levels, except possibly the last have the maximum number of possible nodes, and if all the nodes at the last level appear as far left as possible.

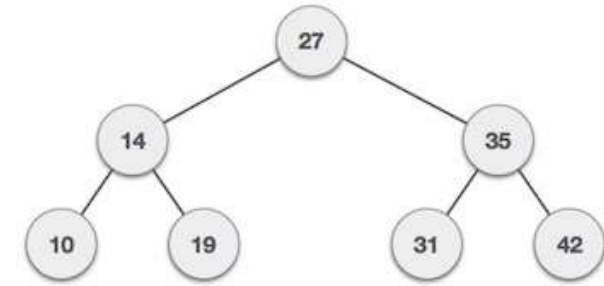


Complete Binary tree

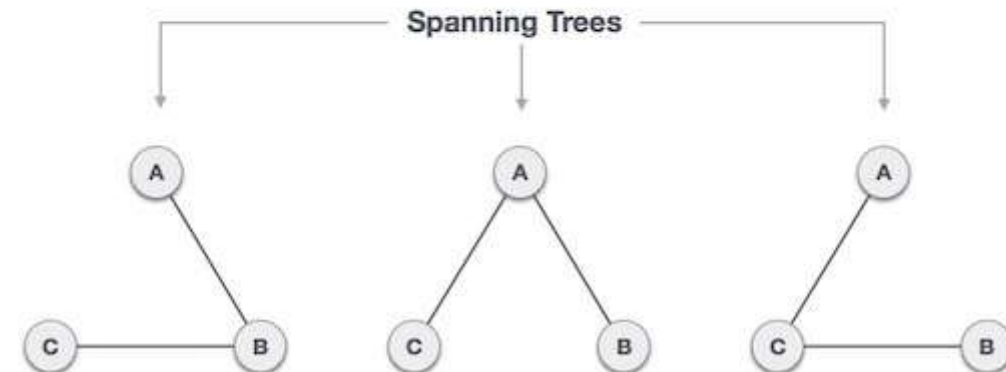
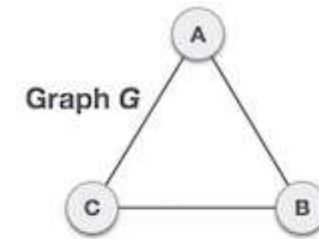
Trees

Types of Trees:

- **Binary Search Tree:** Every node in the left subtree must contain a value less than the value of the root node, and the value of each node in the right subtree must be bigger than the value of the root node.
- **Spanning Tree:** A spanning tree is a subset of Graph G, which has all the vertices covered with minimum possible number of edges. Hence, a spanning tree does not have cycles and it cannot be disconnected..



Binary search tree



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